# How Much the Rounding Errors Could Affect the Computer Results 

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#### Abstract

The paper emphasizes the specific manner to perform the arithmetic operations in the computer as always it is possible to have "rounding errors". Neglecting the influence of the computer rounding errors could affect a right conclusion, for example, in the case of an account balance verification.


Key words : rounding errors, Monte Carlo technique, specific distributions.
JEL Classification System : C87, C15, C16.

## 1. The Problem Formulation

The numerical values are approximated by the computer by retaining only a finite number $q$ of significant digits. For this reason the results of the computer arithmetic operations are rounded to the first $q$ important figures. For example, the Excel spreadsheet operates in the arithmetic calculus with $q=15$ significant figures.
We can obtain wrong results if we don't take into consideration the possibility to have errors in the rounding process of computer outcome. We intend to give a quantitative image regarding the influence of the computer rounding errors.
For the simplicity of the exposure we'll treat the ordinary case of solving a linear system which has only two equations, the variables $x, y$ being unknown, that is

$$
\left\{\begin{array}{l}
a_{1} x+a_{2} y=a_{3}  \tag{1}\\
a_{4} x+a_{5} y=a_{6}
\end{array}\right.
$$

The model (1) is characterized by the input parameters $a_{i}, 1 \leq i \leq 6$.
Obviously, the solution $\left(x_{0}, y_{0}\right)$ of the system (1) has the form:

[^0]\[

$$
\begin{equation*}
x_{0}=\frac{a_{5} a_{3}-a_{2} a_{6}}{a_{1} a_{5}-a_{2} a_{4}} \quad y_{0}=\frac{a_{1} a_{6}-a_{4} a_{3}}{a_{1} a_{5}-a_{2} a_{4}} \tag{2}
\end{equation*}
$$

\]

The presence of computer rounding errors could induce an unexpected effect, that is the "solution" ( $x_{0}, y_{0}$ ) obtained with formulas (2) don't necessary verify both equations (1).
We'll define the indicator $W$ to determine the frequency of this "surprising" non concordance aspect . More precisely, the variable $W$ takes the values 0,1 or 2 when, respectively, the quantities $x_{0}, y_{0}$ satisfy the both equations (1), verify only one of the relations (1), or none of the equalities (1) are fulfilled.
In the following we'll estimate the probability that the situations $W=k, k \in(0,1,2\}$ appear.

## 2. The Distribution of Input Data

We intend to determine the frequencies of the events $W=k, k \in\{0,1,2\}$, when $n$ systems of the type (1) were taken into consideration, the values of the input parameters $a_{i}, 1 \leq i \leq 6$, being chosen randomly.
So, we'll suppose that the quantities $a_{i}$ are observations of the random variables (r.v.) $A_{i}, 1 \leq i \leq 6$, the random vector $\left(A_{1}, A_{2}, A_{3}, A_{4}, A_{5}, A_{6}\right)$ having a known cumulative distribution function ( c.d.f. ) $F_{A}\left(a_{1}, a_{2}, a_{3}, a_{4}, a_{5}, a_{6}\right)$.
Below the random variables $A_{i}$ will have a power type distribution $\operatorname{Pow}(\lambda)$ or also a standard normal distribution $\operatorname{Nor}(0,1)$.
Definition 1. The r.v. $Z$ has a power distribution $\operatorname{Pow}(\lambda), Z \sim \operatorname{Pow}(\lambda), \lambda>0$, if its probability density function ( p.d.f. ) $f_{1}(z ; \lambda)$ is given by the expression:

$$
\begin{equation*}
f_{1}(z ; \lambda)=\lambda z^{\lambda-1}, \quad 0 \leq z \leq 1 \tag{3}
\end{equation*}
$$

The graphic $\mathbf{G 1}$ presents the variation in the p.d.f. $f_{1}(z ; \lambda), 0 \leq z \leq 1$, for different values of the shape power distribution parameter $\lambda \geq 1(\lambda \geq 1)$.

Graphic G1
The p.d.f. $\mathrm{F}_{1}(\mathbf{z} ; \lambda)$

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Remark 1. Interpreting the graphic $\mathbf{G 1}$ we deduce that the r. v. $Z \sim \operatorname{Pow}(\lambda)$ with $\lambda>1$ generates more frequently relative higher $z$ random values.

In general we'll consider that the random variables $A_{i}, 1 \leq i \leq 6$, are identically distributed and independent too. But we'll also treat the situation when two arbitrary input random variables, $A_{i}$ and $A_{j}$, are correlated, that is $\operatorname{Cor}\left(A_{i}, A_{j}\right)=\rho$ with $-1 \leq \rho \leq 1$.
Definition 2. We designate by $(Z, V) \sim \operatorname{Nor2}(\rho)$ if the random vector $(Z, V)$ has a bivariate normal distribution with $Z \sim \operatorname{Nor}(0,1), V \sim \operatorname{Nor}(0,1)$ and $\operatorname{Cor}(Z, V)=\rho$, $-1 \leq \rho \leq 1$.

In the subsequent we'll estimate the probabilities $\operatorname{Pr}(W=k), k \in\{0,1,2\}$, considering the following variants :
Example 1. The input random variables $A_{i}, 1 \leq i \leq 6$, of model (1) are independent and identically distributed with $A_{i} \sim \operatorname{Pow}(\lambda), \lambda>0$.

Example 2. The random vectors $\left(A_{1}, A_{4}\right),\left(A_{2}, A_{5}\right),\left(A_{3}, A_{6}\right)$ are independent and in addition $\left(A_{1}, A_{4}\right) \sim \operatorname{Nor2}(\rho),\left(A_{2}, A_{5}\right) \sim \operatorname{Nor2}(\rho),-1 \leq \rho \leq 1,\left(A_{3}, A_{6}\right) \sim$ Nor2(0).
If $\rho=0$ than all the random variables $A_{i}, 1 \leq i \leq 6$, are independent and have a standard normal distribution $\operatorname{Nor}(0,1)$.

## 3. A Monte Carlo Simulation

The difficulty of the theoretical evaluation of the probabilities $\operatorname{Pr}(W=k)$, $k \in\{0,1,2\}$, arises from the relative complexity of the formulas (2) and also from the sophisticated form of the multivariate distribution $F_{A}\left(a_{1}, a_{2}, \ldots, a_{6}\right)$ which
characterizes the random vector $\left(A_{1}, A_{2}, \ldots, A_{6}\right)$. Moreover, we'll impose a fixed number $m$ of decimal figures for the input parameters $a_{i}$ of the model (1).
For this reason we'll avoid a theoretical approach, preferring the computer Monte Carlo simulations. The stochastic simulation algorithm ASim determines the frequencies $S_{k}$ to create the situations $W=k, k \in\{0,1,2\}$, when $n$ linear systems of type (1) are solved.
Algorithm ASim (computes the frequencies $S_{k}$ for appearing of the events $W=k$ ).
Step 0. Inputs : The number $n$ of the type (1) systems used in the evaluation process.
C.d.f. $F_{A}\left(a_{1}, a_{2}, \ldots, a_{6}\right)$ for the random vector $\left(A_{1}, A_{2}, \ldots, A_{6}\right)$.

The number $m$ of the decimal places imposed for the parameters $a_{i}, 1 \leq i \leq 6$ :
Step 1. $t=0 \quad(t$ counts the number of systems (1) which are solved)

$$
S_{0}=0 \quad S_{1}=0 \quad S_{2}=0
$$

Step 2. $t=t+1$
If $t>n$ then Print $S_{0}, S_{1}, S_{2}$. STOP
Step 3. Generate a random vector $\left(a_{1}, a_{2}, \ldots, a_{6}\right)$ having the c.d.f. $F_{A}\left(a_{1}, a_{2}, \ldots, a_{6}\right)$

Step 4. Determine the approximations $a_{i}^{*}$ (with $m$ decimal places) for $a_{i}$ values, $1 \leq i \leq 6$

Step 5. $x_{0}=\frac{a_{5}^{*} a_{3}^{*}-a_{2}^{*} a_{6}^{*}}{a_{1}^{*} a_{5}^{*}-a_{2}^{*} a_{4}^{*}} \quad y_{0}=\frac{a_{1}^{*} a_{6}^{*}-a_{4}^{*} a_{3}^{*}}{a_{1}^{*} a_{5}^{*}-a_{2}^{*} a_{4}^{*}}$
Step 6. $W=0$
If $a_{1}^{*} x+a_{2}^{*} y=a_{3}^{*}$ then $W=W+1$
If $a_{4}^{*} x+a_{5}^{*} y=a_{6}^{*}$ then $W=W+1$
Step 7. If $W=0$ then $S_{0}=S_{0}+1$
If $W=1$ then $S_{1}=S_{1}+1$
If $W=2$ then $S_{2}=S_{2}+1$
Go to Step 2.
Remark 2. The simulation algorithm ASim was implemented in Excel spreadsheet. To obtain observations from a uniform $[0,1]$ distributed r.v. $U$, $U \sim \operatorname{Uni}([0,1])$, we run the Excel procedure Rand. The generation of the random values $Z, Z \sim$ $\operatorname{Uni}([0,1])$, is based on the inverse method (Gentle [3], p. 42-43 ).

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Applying the following proposition we reduce the generation of the random values $Z$, $Z \sim \operatorname{Pow}(\lambda), \lambda>0$, by producing observations from $U \sim \operatorname{Uni}([0,1])$.
Proposition 1. If the r.v. $U$ has a uniform $[0,1]$ distribution and $Z=U^{1 / \lambda}, \lambda>0$, then $Z \sim \operatorname{Pow}(\lambda)$.
Proof. We'll compute the c.d.f. $F_{1}(z ; \lambda)$ of the r.v. $Z=U^{1 / \lambda}$ for an arbitrary $0 \leq z \leq 1$. So,

$$
F_{1}(z)=\operatorname{Pr}(Z \leq z)=\operatorname{Pr}\left(U^{1 / \lambda} \leq z\right)=\operatorname{Pr}\left(U \leq z^{\lambda}\right)=z^{\lambda}
$$

Therefore the p.d.f. $f_{1}(z ; \lambda)$ of the r.v. $Z$ is given by the expression

$$
f_{1}(z ; \lambda)=\frac{\partial F_{1}(z ; \lambda)}{\partial z}=\lambda z^{\lambda-1}
$$

that is $Z \sim \operatorname{Pow}(\lambda)$.
The generation of the random vector $(Z, V),(Z, V) \sim \operatorname{Nor} 2(\rho),-1 \leq \rho \leq 1$, is based on the following result :
Proposition 2. Let $F_{0}(z), x \in R$, be the c.d.f. of the r.v. $Z, Z \sim \operatorname{Nor}(0,1)$ (Laplace function). If $-1 \leq \rho \leq 1, U_{1}, U_{2}$ are independent random variables having a uniform [ 0,1 ] distribution and

$$
\begin{gather*}
Z=F_{0}^{-1}\left(U_{1}\right) \\
V=\rho Z+\sqrt{1-\rho^{2}} F_{0}^{-1}\left(U_{2}\right) \tag{4}
\end{gather*}
$$

then $(Z, V) \sim \operatorname{Nor} 2(\rho)$.
Proof. Obviously, the r.v.-s $Z=F_{0}^{-1}\left(U_{1}\right)$ and $F_{0}^{-1}\left(U_{2}\right)$ are independent and identically $\operatorname{Nor}(0,1)$ distributed ( Gentle [3], p.42-43, the inverse method). In this conditions the r.v. $V$ defined by formula (4) is $\operatorname{Nor}(0,1)$ distributed (Papoulis [6]) and, in addition, the random vector $(Z, V)$ has a bivariate normal distribution (Papoulis [6], p.162-164 ; Gentle [3], p.105-109). More,

$$
\begin{gathered}
\operatorname{Cor}(Z, V)=\frac{\operatorname{Cov}(Z, V)}{\sqrt{\operatorname{Var}(Z)} \sqrt{\operatorname{Var}(V)}}=\frac{\operatorname{Cov}\left(Z, \rho Z+\sqrt{1-\rho^{2}} F_{0}^{-1}\left(U_{2}\right)\right)}{\sqrt{\rho^{2} \operatorname{Var}(Z)+\left(1-\rho^{2}\right) \operatorname{Var}\left(F_{0}^{-1}\left(U_{2}\right)\right)}}= \\
=\frac{\rho \operatorname{Cov}(Z, Z)+\sqrt{1-\rho^{2}} \operatorname{Cov}\left(Z, F_{0}^{-1}\left(U_{2}\right)\right)}{\sqrt{\rho^{2}+\left(1-\rho^{2}\right)}}=\rho \operatorname{Var}(Z)=\rho
\end{gathered}
$$

which prove that $(Z, V) \sim \operatorname{Nor} 2(\rho)$.
To increase the accuracy of Monte Carlo estimations, the algorithm ASim will be run successively many times.
Tables 1a, 1b, 1c present the resulted frequencies $S_{k}$ that the events $W=k$, $k \in\{0,1,2\}$ appear, where $S_{0}+S_{1}+S_{2}=n=1000$. The input parameters $a_{i}$ of ASim simulation procedure respects one of the restrictions mentioned in Examples 1 or 2 . Constantly, the simulation algorithm ASim was applied 10 times, the index $S$ designating the current simulation execution, $1 \leq s \leq 10$.
The estimation of the probabilities $\operatorname{Pr}(W=k), \quad k \in\{0,1,2\}$, taking into consideration the average results of all 10 simulations, are given in Tables 2a, 2b, 2c. In fact Tables $2 a, 2 b, 2 c$ convert in percent the mean of the same frequency $S_{k}$ for all simulation values listed in the similar Tables $2 a, 2 b$, respectively $2 c$.

Table 1a
The frequency of the event $W=k, k \in\{0,1,2\}$, at the simulation $s$
(Example 1, $n=1000, m=20$ )

| $\lambda$ | W | $\mathrm{s}=1$ | $\mathrm{s}=2$ | s=3 | $\mathrm{s}=4$ | s=5 | $\mathrm{s}=6$ | s=7 | s=8 | s=9 | s=10 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1/5 | 0 | 576 | 563 | 563 | 591 | 589 | 598 | 629 | 549 | 549 | 582 |
|  | 1 | 419 | 424 | 426 | 398 | 403 | 391 | 369 | 441 | 438 | 408 |
|  | 2 | 5 | 13 | 11 | 11 | 8 | 11 | 2 | 10 | 13 | 10 |
| 1/4 | 0 | 619 | 588 | 622 | 647 | 610 | 614 | 613 | 640 | 647 | 628 |
|  | 1 | 368 | 403 | 362 | 337 | 384 | 374 | 377 | 349 | 344 | 359 |
|  | 2 | 13 | 9 | 16 | 16 | 6 | 12 | 10 | 11 | 9 | 13 |
| 1/3 | 0 | 678 | 692 | 651 | 652 | 672 | 689 | 669 | 679 | 682 | 677 |
|  | 1 | 307 | 293 | 329 | 331 | 305 | 290 | 319 | 309 | 308 | 307 |
|  | 2 | 15 | 15 | 20 | 17 | 23 | 21 | 12 | 12 | 10 | 16 |
| 1/2 | 0 | 761 | 750 | 720 | 759 | 791 | 758 | 760 | 776 | 734 | 743 |
|  | 1 | 213 | 217 | 254 | 212 | 181 | 213 | 218 | 202 | 249 | 233 |
|  | 2 | 26 | 33 | 26 | 29 | 28 | 29 | 22 | 22 | 17 | 24 |
| 1 | 0 | 863 | 838 | 821 | 848 | 843 | 853 | 840 | 856 | 826 | 843 |
|  | 1 | 94 | 106 | 125 | 104 | 88 | 102 | 105 | 103 | 114 | 102 |
|  | 2 | 43 | 56 | 54 | 48 | 69 | 45 | 55 | 41 | 60 | 55 |
| 2 | 0 | 806 | 821 | 812 | 811 | 814 | 826 | 826 | 810 | 809 | 810 |
|  | 1 | 64 | 73 | 81 | 75 | 80 | 66 | 73 | 63 | 74 | 75 |
|  | 2 | 130 | 106 | 107 | 114 | 106 | 108 | 101 | 127 | 117 | 115 |
| 3 | 0 | 746 | 735 | 742 | 751 | 741 | 761 | 742 | 739 | 753 | 764 |
|  | 1 | 60 | 68 | 71 | 68 | 84 | 67 | 73 | 76 | 66 | 74 |
|  | 2 | 194 | 197 | 187 | 181 | 175 | 172 | 185 | 185 | 181 | 162 |
| 4 | 0 | 690 | 677 | 698 | 671 | 665 | 676 | 711 | 696 | 653 | 699 |
|  | 1 | 74 | 65 | 84 | 76 | 78 | 78 | 60 | 71 | 87 | 68 |

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|  | 2 | 236 | 258 | 218 | 253 | 257 | 246 | 229 | 233 | 260 | 233 |
| ---: | :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| 5 | 0 | 654 | 615 | 611 | 653 | 642 | 630 | 636 | 625 | 647 | 644 |
|  | 1 | 70 | 78 | 79 | 82 | 77 | 71 | 70 | 60 | 76 | 87 |
|  | 276 | 307 | 310 | 265 | 281 | 299 | 294 | 315 | 277 | 269 |  |

Analyzing the values from the mentioned tables we conclude :

- The frequencies $S_{0}, S_{1}, S_{2}$ for the events $W=0, W=1, W=2$ are very distinct and depend on the specific form of the c.d.f. $F_{A}\left(a_{1}, a_{2}, \ldots, a_{6}\right)$ which characterizes the parameters $a_{i}$ of the model (1) (Tables 1a, 2a, 1b, 2b).
Moreover, we also remark important variations of the frequency $S_{0}$ when we change the distribution of the random variables $A_{i}, 1 \leq i \leq 6$, but preserving the independence of these variables (Tables 1a, 1b; Tables 2a, $2 b$ - case $\rho=0$ ).

Table 1b
The estimation of $\operatorname{Pr}(W=k), k \in\{0,1,2\}$, from 10 simulations
(Example 1, $n=1000, m=20$ )

| W | $\lambda=1 / 5$ | $\lambda=1 / 4$ | $\lambda=1 / 3$ | $\lambda=1 / 2$ | $\lambda=1$ | $\lambda=2$ | $\lambda=3$ | $\lambda=4$ | $\lambda=5$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0.5789 | 0.6228 | 0.6741 | 0.7552 | 0.8431 | 0.8145 | 0.7474 | 0.6836 | 0.6357 |
| 1 | 0.4117 | 0.3657 | 0.3098 | 0.2192 | 0.1043 | 0.0724 | 0.0707 | 0.0741 | 0.0750 |
| 2 | 0.0094 | 0.0115 | 0.0161 | 0.0256 | 0.0526 | 0.1131 | 0.1819 | 0.2423 | 0.2893 |

Table 2a
The frequency of the event $W=k, k \in\{0,1,2\}$, at the simulation $s$
(Example 2, $n=1000, m=20$ )

| $\rho$ | W | $\mathrm{s}=1$ | $\mathrm{s}=2$ | s=3 | $\mathrm{s}=4$ | s=5 | s=6 | s=7 | $\mathrm{s}=8$ | $\mathrm{s}=9$ | $\mathrm{s}=10$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0.0 | 0 | 801 | 822 | 804 | 810 | 817 | 821 | 814 | 814 | 836 | 808 |
|  | 1 | 180 | 163 | 177 | 174 | 163 | 163 | 167 | 176 | 151 | 177 |
|  | 2 | 19 | 15 | 19 | 16 | 20 | 16 | 19 | 10 | 13 | 15 |
| 0.4 | 0 | 790 | 809 | 785 | 787 | 810 | 820 | 814 | 813 | 817 | 788 |
|  | 1 | 183 | 169 | 193 | 194 | 163 | 156 | 173 | 172 | 168 | 196 |
|  | 2 | 27 | 22 | 22 | 19 | 27 | 24 | 13 | 15 | 15 | 16 |
| 0.8 | 0 | 769 | 737 | 750 | 731 | 768 | 780 | 765 | 758 | 748 | 776 |
|  | 1 | 206 | 229 | 211 | 219 | 205 | 189 | 193 | 212 | 216 | 185 |
|  | 2 | 25 | 34 | 39 | 50 | 27 | 31 | 42 | 30 | 36 | 39 |
| 0.9 | 0 | 718 | 695 | 738 | 703 | 737 | 730 | 721 | 713 | 713 | 705 |
|  | 1 | 230 | 255 | 223 | 245 | 214 | 225 | 232 | 238 | 243 | 246 |
|  | 2 | 52 | 50 | 39 | 52 | 49 | 45 | 47 | 49 | 44 | 49 |
| 0.95 | 0 | 696 | 681 | 689 | 684 | 687 | 684 | 702 | 670 | 635 | 651 |
|  | 1 | 238 | 261 | 249 | 246 | 250 | 258 | 236 | 265 | 292 | 282 |


| 2 | 66 | 58 | 62 | 70 | 63 | 58 | 62 | 65 | 73 | 67 |
| :--- | :--- | :--- | :--- | :--- | ---: | ---: | ---: | ---: | ---: | ---: |

- An imposed strong correlation coefficient $\rho$ between the parameters $a_{i}$, $1 \leq i \leq 6$, affects essentially the frequencies $S_{0}, S_{1}, S_{2}$ of the events $W=0, W=1, W=2($ Tables $2 a$ and $2 b)$.
- Contrary, the use of the random inputs $a_{i}, 1 \leq i \leq 6$, which have an imposed number $m$ of decimal places do not often affect much the frequencies $S_{k}$, $k \in\{0,1,2\}$ (Tables 3a, 3b). This aspect could be explained by the fact that the situations $W=k, k \in\{0,1,2\}$, depend basically by the complexity of the formulas (2), being secondary affected by the accuracy of the approximations $a_{i}^{*}$ of the values $a_{i}, 1 \leq i \leq 6$.

Table 2b
The estimation of $\operatorname{Pr}(W=k), k \in\{0,1,2\}$, from 10 simulations
(Example 2, $n=1000, m=20$ )

|  | $\rho=0.0$ | $\rho=0.4$ | $\rho=0.8$ | $\rho=0.9$ | $\rho=0.95$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $W=0$ | 0.8147 | 0.8033 | 0.7582 | 0.7173 | 0.6779 |
| $W=1$ | 0.1691 | 0.1767 | 0.2065 | 0.2351 | 0.2577 |
| $W=2$ | 0.0162 | 0.0200 | 0.0353 | 0.0476 | 0.0644 |

Table 3a
The frequency of the event $W=k, k \in\{0,1,2\}$, at the simulation $s$
(Example 1, $n=1000, \lambda=1$ )

| m | W | $\mathrm{s}=1$ | $\mathrm{s}=2$ | $\mathrm{s}=3$ | $\mathrm{s}=4$ | $\mathrm{s}=5$ | $\mathrm{s}=6$ | $\mathrm{s}=7$ | $\mathrm{s}=8$ | $\mathrm{s}=9$ | $\mathrm{s}=10$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 12 | 0 | 838 | 812 | 824 | 819 | 858 | 822 | 810 | 834 | 834 | 833 |
|  | 1 | 111 | 124 | 120 | 122 | 99 | 125 | 135 | 109 | 102 | 114 |
|  | 2 | 51 | 64 | 56 | 59 | 43 | 53 | 55 | 57 | 64 | 53 |
| 13 | 0 | 819 | 825 | 830 | 830 | 837 | 814 | 825 | 824 | 835 | 823 |
|  | 1 | 135 | 125 | 118 | 114 | 112 | 121 | 130 | 114 | 124 | 131 |
|  | 2 | 46 | 50 | 52 | 56 | 51 | 65 | 45 | 62 | 41 | 46 |
| 14 | 0 | 828 | 831 | 829 | 844 | 825 | 832 | 839 | 834 | 836 | 834 |
|  | 1 | 119 | 114 | 120 | 106 | 115 | 114 | 117 | 115 | 114 | 118 |
|  | 2 | 53 | 55 | 51 | 50 | 60 | 54 | 44 | 51 | 50 | 48 |
| 15 | 0 | 835 | 821 | 843 | 820 | 821 | 841 | 845 | 821 | 822 | 838 |
|  | 1 | 118 | 128 | 106 | 122 | 131 | 115 | 106 | 125 | 123 | 114 |
|  | 2 | 47 | 51 | 51 | 58 | 48 | 44 | 49 | 54 | 55 | 48 |
| 20 | 0 | 843 | 863 | 831 | 864 | 845 | 853 | 860 | 844 | 839 | 839 |
|  | 1 | 109 | 93 | 113 | 97 | 102 | 111 | 88 | 103 | 104 | 111 |


| 2 | 48 | 44 | 56 | 39 | 53 | 36 | 52 | 53 | 57 | 50 |
| :--- | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |

Table 3b
The estimation of $\operatorname{Pr}(W=k), k \in\{0,1,2\}$, from 10 simulations
(Example 1, $n=1000, \lambda=1$ )

| $W$ | $\mathrm{~m}=12$ | $\mathrm{~m}=13$ | $\mathrm{~m}=14$ | $\mathrm{~m}=15$ | $\mathrm{~m}=20$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0.8284 | 0.8262 | 0.8332 | 0.8307 | 0.8481 |
| 1 | 0.1161 | 0.1224 | 0.1152 | 0.1188 | 0.1031 |
| 2 | 0.0555 | 0.0514 | 0.0516 | 0.0505 | 0.0488 |

- Graphic G2 presents suggestively the variation in the probabilities $\operatorname{Pr}(W=0)$ depending on the shape parameter $\lambda$ which defines the power distribution for the random values $a_{i}, 1 \leq i \leq 6$ ( Table $1 b$ ). Interpreting simultaneously the graphics G1 and G2 we remark a "monotony relation" between the probabilities $\operatorname{Pr}(W=0)$ and the $\lambda$ values $(0<\lambda \leq 1$ or $\lambda \geq 1$; see also Remark 1$)$.

Graphic G2
The estimation of $\operatorname{Pr}(\mathbf{W}=\mathbf{0})$


## 4. Concluding Remarks

The computer process which determines the solution $\left(x_{0}, y_{0}\right)$ of the system (1) is affected by the rounding errors. The influence of the computer approximation errors is emphasized by the non-zero frequencies that the events $W=1$ or $W=2$ appear.

Surprising, but perfectly explained, is the fact that the probability $\operatorname{Pr}(W=0)$, to solve "correctly" with the computer the system (1), is not closed to the maximal value 1. Tables 1a, 1b, 1 c mention more situations where $\operatorname{Pr}(W=0) \approx 0.5$.

Obviously, the probability $\operatorname{Pr}(W=0)$ depends on the distribution $F_{A}\left(a_{1}, a_{2}, \ldots, a_{6}\right)$ of the parameters $a_{1}, a_{2}, \ldots, a_{6}$ which defines system (1). So, we studied the situation when the r.v. $A_{i}$ were independent, being analyzed the correlation aspect too. The tables $1 a-1 c, 2 a-2 c$ and the graphs G1-G2 give us a quantitative image regarding the influence of computer rounding errors.
The present paper warns about the possibility to produce errors even if we use a computer. Neglecting in practice the influence of the computer rounding errors could imply a lot of wrong decisions (see, for example, the computational errors in accounting, [8]).

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    (BR) - Romanian Journal of Economic Forecasting - 2/2007

